

We work over a field \mathbb{K} . $\text{pt} := \text{Spec } \mathbb{K}$.

Def. An alg. stack Y of finite type is pointy if $Y \xrightarrow{\text{open, dense}} \text{pt}$.

(If you wish, density means "schematically dense" rather than "topologically".)

Example. $Y = X/G$, $X \supset G$ $\xrightarrow{\text{open dense}}$. E.g., $Y = \mathbb{A}^1/\mathbb{G}_m$.

Now let C be a smooth curve (e.g., \mathbb{A}^1).

Stack $\text{Maps}^\circ(C, Y)$ (stack of nondegenerate maps $C \rightarrow Y$):

an S -point of $\text{Maps}^\circ(C, Y)$ is $f: C \times S \rightarrow Y$ s.t. $\bar{f}(y|_{\text{pt}})$ is finite over S .

In the future I will describe \mathbb{K} -points of stacks and pretend that this is enough to understand S -points. in practice it is over S .

Pre-theorem. Assume $Y \xrightarrow{\text{diag}} Y \times Y$ separated (rather than quasi-separated). affine

Then $\text{Maps}^\circ(C, Y)$ is an alg. space locally of finite type. In some important cases this is a theorem, with $\text{Maps}^\circ(C, Y)$ being a scheme.

Exercise. $\text{Maps}^\circ(C, Y)$ is a "space", i.e., S -points have no (nontrivial) automorphisms. (It is here that the separatedness assumption fine bundle on \mathcal{D} is used.)

Example. $\text{Maps}^\circ(C, \mathbb{A}^1/\mathbb{G}_m) = \{(L, s)\} \mid L \in H^0(C, \mathcal{O}) \neq 0\} = \text{Div}(C) =$
 $= \text{Sym}^n C := \coprod_{n \geq 0} \text{Sym}^n C$, the scheme of zeros of s is finite

Exercise. $\text{Maps}^\circ(C, \mathbb{P}^1/\mathbb{G}_a) = ?$ [Answer: $\{(D, \sigma) \mid D \subset C, \sigma \in H^0(D, \mathcal{O}_D)\}$]

(Is it true that $\text{Maps}(\text{Spec } \mathbb{K}[E_8], \mathbb{A}^1/\mathbb{G}_a) = \mathbb{P}^1/\mathbb{G}_a$? " tangent bundle of $\text{Sym}^n C$. Note: $\mathbb{P}^1 \not\simeq \mathbb{P}^1/\mathbb{G}_a = \mathbb{P}^1 \otimes \mathbb{PGL}(2)/N$. Here is a generalization:

Example. G reductive, $X = G/N$, $\text{Maps}^\circ(C, Y) = B \times^G X$, $Y = B \times^G N$ is pointy.

Then $\text{Maps}^\circ(C, Y)$ known to be a scheme; this is Mirkovic + Finkelberg's scheme.

In this example the complement of the open orbit is a Cartier divisor.

Generalization. H affine alg. group, $Y = X/H$ (where X is an H -variety).

Assume $X \supset H$. Assume \exists H -invariant Cartier divisor $\Delta \subset X$ such that

$X \setminus \Delta = H$. We'll prove the quasi-finite theorem if X is quasi-affine. We'll construct a morphism from $\text{Maps}^\circ(C, Y)$ to a certain scheme.

Have $\text{Maps}^\circ(C, Y) \xrightarrow{\psi} \text{Sym}^n C$. Have something better:

$$\psi \mapsto \mathcal{D}_f := \bar{f}(\Delta/H)$$

We have $\text{Maps}^\circ(C, Y) \xrightarrow{\psi} \text{Maps}^\circ(C, \text{pt}/H) = \{H\text{-bundles on } C\}$

$$\psi \mapsto \mathcal{E}_f \longrightarrow \mathcal{E}_f$$

\mathcal{E}_f is trivialized over $C \setminus \mathcal{D}_f$. (is an

$\mathcal{G}\mathcal{R}_{\mathcal{D}} := \{H\text{-bundles on } C \text{ trivialized over } C - \mathcal{D}\}$ (this is a sheaf)

$\mathcal{G}\mathcal{R} \rightarrow \text{Sym}^n C$, $\mathcal{G}\mathcal{R}_{\mathcal{D}} = \text{fiber over } \mathcal{D}$.

Thus we get $\text{Maps}^\circ(C, Y) \rightarrow \text{GR}$. not merely ind-scheme!

Prove (or maybe disprove) that: $\text{Maps}^\circ(C, Y) \rightarrow \text{GR}$.

Exercise. 1) X affine $\Rightarrow \text{Maps}^\circ(C, Y)$ is a closed subscheme of GR

X quasi-affine \Rightarrow locally closed

2) $\mathcal{O}_X(A)$ ample $\Rightarrow \text{Maps}^\circ(C, Y)$ is a locally closed subscheme of GR

This assumption holds if X is quasi-affine.

Relation to formal arcs.

Grinberg-Kazhdan. $X \supset X^{sm}$ - smooth open locus

finite type

(γ describes movement of a point on X)

$\gamma \in (\mathcal{L}X)(k)$ $\gamma: \text{Spec } k[[t]] \rightarrow X$, $t = \text{"time"}, \gamma(t) \in X^{sm}$ for $t \neq 0$.

$\mathcal{L}X := \text{Maps}(D, X)$, $D := \text{Spec } k[[t]]$, $\mathcal{L}X$ is a scheme of infinite type

$\gamma \in (\mathcal{L}X)(k)$.

Thm. $\mathcal{L}X_y \cong \mathcal{Y}_y \times D^*$ for some pair (Y, y) , where Y is some scheme of finite type and $y \in Y(k)$.

(Before Grinberg-Kazhdan), Finkelberg-Mirkovic proved for $X = \overline{G/N}$; they showed that one can take $Y = \text{Maps}^\circ(A^1, X/B)$.

Summary of their argument. (in a slightly more general situation) (Zastava space).

Suppose that 1) an affine alg. group H acts on X , $X \supset H$,

2) $\gamma(t) \in H$ for $t \neq 0$,

3) $\text{Maps}^\circ(A^1, X/H)$ is a scheme of locally finite type (i.e., the quasi-theorem holds)

Then Grinberg-Kazhdan holds for $Y = \text{Maps}^\circ(A^1, X/H)$,

(If $X = \overline{G/N}$ take $H = gBg^{-1}$ for a suitable $g \in G$).

What is $y \in Y(k)$? $y: A^1 \rightarrow X/H$
 \cup
 \cup
 $G_m \rightarrow pt$

Have a covering $G_m \amalg \text{Spec } k[[t]] \rightarrow A^1$ ("morally", an open covering).

Specifying y amounts to specifying $\hat{y}: \text{Spec } k[[t]] \rightarrow X/H$ such that $\text{Spec } k[[t]]$ goes to $pt \subset X/H$ (Believable, can be proved).

Take $\hat{y} = \bar{\gamma}: \text{Spec } k[[t]] \rightarrow X/H$,

follows from Beauville-Baszle

(-3-)

How does one get an isomorphism $\widehat{LX}_g \xrightarrow{\sim} \widehat{Y}_y \times D^\infty$?

There ~~is~~ It is not quite canonical, canonically, one has $\widehat{LX}_g \rightarrow \widehat{Y}_y$, torsor over ^{certain} formal group scheme g .

As a formal scheme, $g \cong D^\infty$. Choosing such an isomorphism and also a section of the torsor, get Grinberg-Kazhdan.

Description of g : H acts on X , LH acts on LX , so

$g := \widehat{LH}$ acts on \widehat{LX}_g . \checkmark As a formal scheme, $H \cong D^\infty$, so $\widehat{LH} \cong D^\infty$.

For ~~any~~ general X , the Finkelberg-Mirkovic method has to be modified.

(E.g., if X is a curve whose normalization has genus > 1 then ~~any~~ connected algebraic group cannot act on X nontrivially.)

Modification: instead of group actions use (smooth) groupoids acting on X . In my article on Grinberg-Kazhdan this idea is used "behind the scenes" (there is only a hint in the last paragraph). I will try to explain the hint.

(4)

Generalities on groupoids.

(abstract) group group scheme

(abstract) groupoid groupoid in {Schemes}

Abstract groupoids,

Def. A groupoid is a category in which all morphisms are invertible.

$$X = \{\text{objects}\}, \quad \Gamma = \{\text{morphisms}\} \quad \left. \begin{array}{l} \Gamma \xrightarrow{P_1} X \quad (\text{source, target}). \\ \Gamma \times_X \Gamma \xrightarrow{c} \Gamma \quad \text{composition.} \end{array} \right\} \text{Data}$$

Certain properties should hold. In particular, one should have the "identity" $\underset{x \in X}{\exists} X \xrightarrow{e} \Gamma$ and inversion $\Gamma \xrightarrow{i} \Gamma$

If you wish, you can consider e and i as data; then all the properties become identities (no existence quantifiers).

Groupoid in {Schemes} (or in any category in which fiber products exist).

Two equivalent ways:

① To formulate the notion of composition and the above-mentioned identities one only needs the notion of fiber product, and we have this notion in {Schemes};

② A lazier way is to use the language of S-points:

Data: schemes $X, \Gamma, \sqrt{\Gamma} \xrightarrow{P_1} X, \Gamma \times_X \Gamma \xrightarrow{c} \Gamma$

Condition: for any scheme S the diagrams

$$\Gamma(S) \xrightarrow{P_1} X(S), \quad \Gamma(S) \times_{X(S)} \Gamma(S) \xrightarrow{c} \Gamma(S)$$

define a groupoid.

Conventions.

1. I will say "groupoid $\Gamma \Rightarrow X$ " (without mentioning composition)
or "groupoid $\Gamma \rightarrow X \times X$ ".

2. One says " Γ groupoid Γ acting on X ".

Examples of groupoids (say, in $\{\text{Schemes}\}$).

(1) An equivalence relation on X is a groupoid $\Gamma \xrightarrow{\quad} X \times X$ such that $\Gamma \xrightarrow{\quad} X \times X$ is a monomorphism.

Example: $X \times X \xrightarrow{\quad} X$ is a groupoid.

(2) Suppose a group G acts on X . Then we have a groupoid Γ acting on X : a morphism $x_1 \rightarrow x_2$ is an element $g \in G$ such that $gx_1 = x_2$. So the groupoid is $G \times X \xrightarrow[p_1]{\quad} X$, p_1 = projection, p_2 = action.

(3) A group scheme over X is the same as a groupoid $\Gamma \xrightarrow[p_1]{\quad} X$ with $p_1 = p_2$.

(4) $\Gamma \xrightarrow{\quad} X$, $U \xrightarrow{f} X$

Pullback groupoid: $\{\text{objects}\} = U$, a morphism $u_1 \rightarrow u_2$ is a morphism $f(u_1) \rightarrow f(u_2)$.

So we get a groupoid $\Gamma_U \xrightarrow{\quad} U$, $p_U = \Gamma \times_{X \times U} (U \times U)$.

E.g., given a groupoid in $\{\text{Schemes}\}$ acting on X , you can restrict it to an open subscheme of X . Even if Γ comes from a group action, Γ_U usually doesn't.

This example (and many others) shows that groupoids are much more flexible than group actions.

(By the way, groupoids are discussed in Ch. IV of SGA3-I).

Smooth groupoids and algebraic stacks.

Let $\Gamma \xrightarrow[p_1]{\quad} X$ be a groupoid in $\{\text{Schemes}\}$.

Def. A groupoid is smooth if p_1 is smooth (equivalently, if p_2 is smooth).

Note: p_1 and p_2 (i.e., source and target) have equal rights because the inversion map $i: \Gamma \rightarrow \Gamma$ interchanges them.

Suppose we have a smooth groupoid Γ on X such that the map $\Gamma \xrightarrow{\quad} X \times X$ is quasi-compact and separated (or quasi-separated, this depends on the book). To the data one associates the quotient stack $\Gamma \backslash X$ (I prefer to skip the details). A stack is said to be algebraic if it can be represented in this way (with Γ being an alg. stack and X a scheme if you wish).

Details.

X/Γ is the stack (for the smooth topology) associated to the pre-stack $S \mapsto \{\text{groupoid } \Gamma(S) \Rightarrow X(S)\}$.

(Here the "substance" is in "sheafification", i.e., in passing from the pre-stack to the associated stack.)

Here is an "explicit" description of $(X/\Gamma)(S)$, which is hard to find in the literature: an S -point of X/Γ is the following data:

(i) a smooth surjective morphism $S' \rightarrow S$, where S' is an algebraic space,

(ii) a morphism of groupoids

$$\begin{array}{ccc} S' \times_S S' & \rightarrow & \Gamma \\ \downarrow & & \downarrow \\ S' & \rightarrow & X \end{array}$$

such that the squares

$$\begin{array}{ccc} S' \times_S S' & \rightarrow & \Gamma \\ p_1 \downarrow & & \downarrow p_1 \\ S' & \rightarrow & X \end{array}$$

$$\begin{array}{ccc} S' \times_S S' & \rightarrow & \Gamma \\ p_2 \downarrow & & \downarrow p_2 \\ S' & \rightarrow & X \end{array}$$

are Cartesian; in fact, it suffices to require one of the squares to be Cartesian.

(Given a morphism $S \rightarrow X/\Gamma$, one sets $S' := S \times_{X/\Gamma} X$.)